

Region automaton R for timed automaton G

• State space $X \times \mathbb{R}_+^{|\mathcal{C}|}$ is infinite

However, only finitely many Nerode equivalence classes, called regions, exist.

• $n_z \in \mathbb{N}$ be the largest integer against which clock z compared

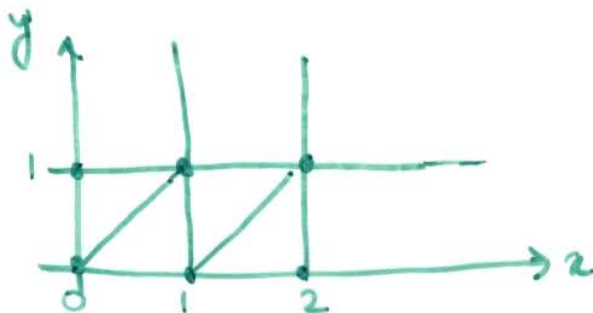
• $\vec{v}_1 \cong \vec{v}_2$ iff

$$(1) \forall z \in \mathcal{C} : \lfloor \vec{v}_1(z) \rfloor = \lfloor \vec{v}_2(z) \rfloor \text{ or } \left[\vec{v}_1(z) > n_z \wedge \vec{v}_2(z) > n_z \right]$$

(2) $\forall z_1, z_2 \in \mathcal{C}$ such that $\vec{v}_i(z_j) \leq n_{z_j}$:

$$\left[\text{frac}(\vec{v}_1(z_1)) \leq \text{frac}(\vec{v}_1(z_2)) \right] \Leftrightarrow \left[\text{frac}(\vec{v}_2(z_1)) \leq \text{frac}(\vec{v}_2(z_2)) \right]$$

Example: $\mathcal{C} = \{x, y\}$; $n_x = 2$, $n_y = 1$



clock space = \mathbb{R}_+^2

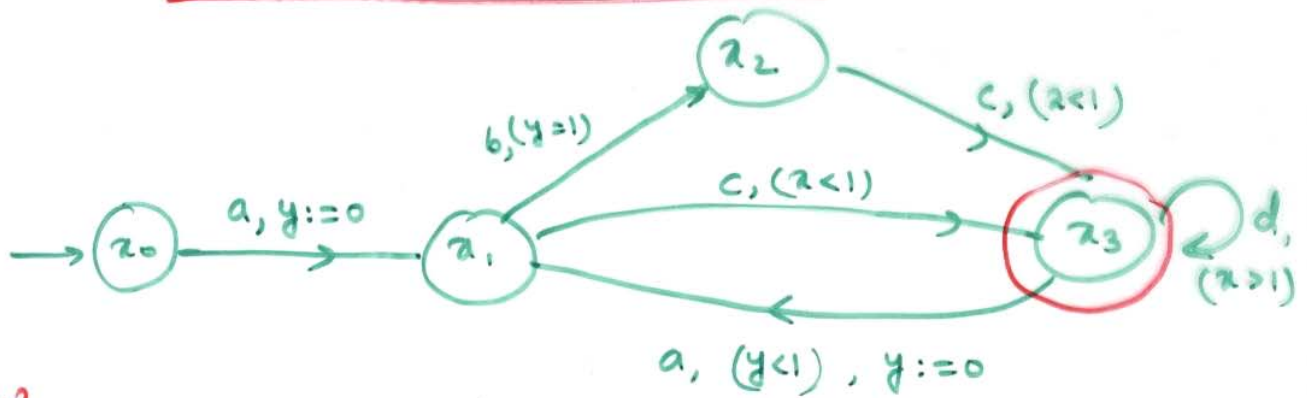
6 corner points : (1,1)

14 open line segments : $0 < x = y < 1$

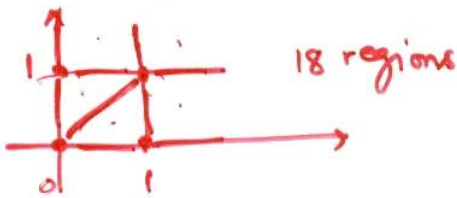
8 open regions : $0 < x < y < 1$

} \Rightarrow 28 clock regions

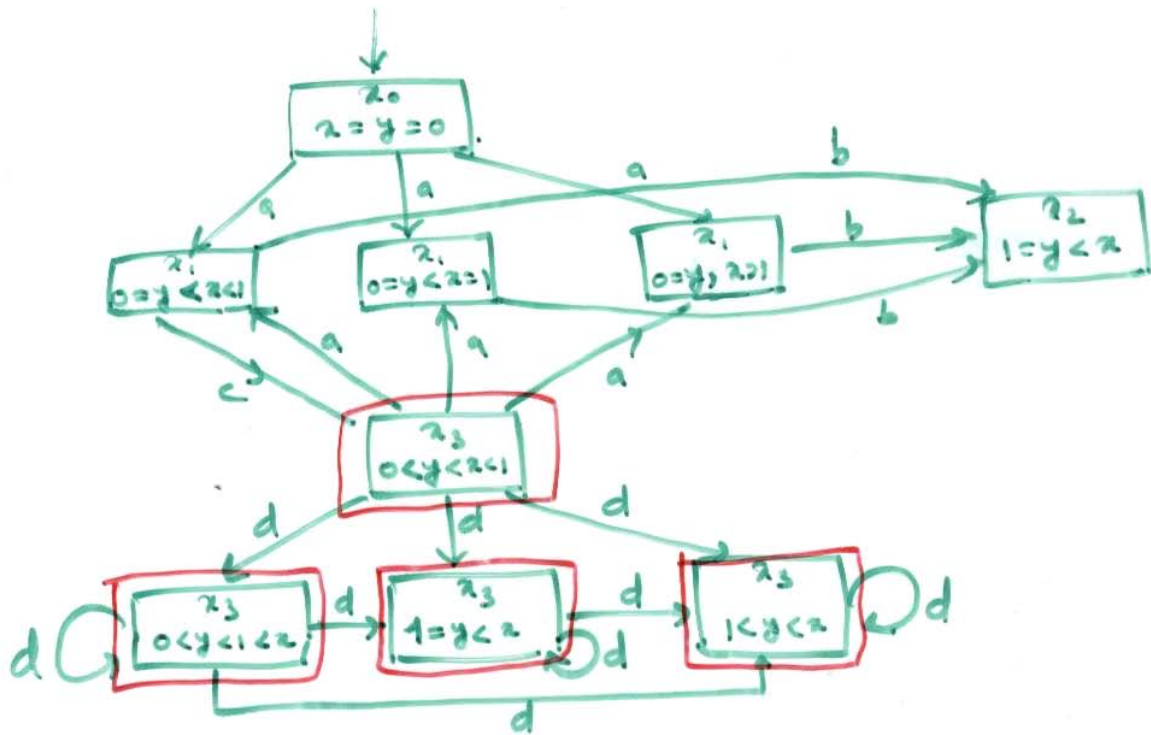
Construction of region automaton R



$C = \{a, y\}$
 $m_a = m_y = 1$



region automaton



$L_m(R) = \text{untime}(T_m(G)) = a(ca)^*cd^*$

Complexity: # of clock regions = $|C|! \cdot 2^{|C|} \prod_{z \in C} (2n_z + 2)$

$\forall z: \{z=c \mid c=0, 1, \dots, n_z\} \cup \{c_1 < z < c_2 \mid c_1, \dots, n_z\} \cup \{z < c_2\} \rightarrow 2n_z + 2$

$\forall z_1, z_2: c_1 < z_1 < c_2, d_1 < z_2 < d_2$ whether $\text{fract}(z_1) \leq \text{fract}(z_2)$

\Rightarrow Emptiness PSPACE complete